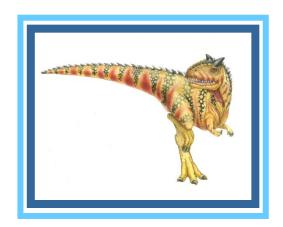
Chapter 13: File-System Interface

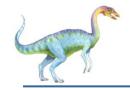




Chapter 13: File-System Interface

- File Concept
- Access Methods
- Disk and Directory Structure
- File-System Mounting
- File Sharing
- Protection





Objectives

- To explain the function of file systems
- To describe the interfaces to file systems
- To discuss file-system design tradeoffs, including access methods, file sharing, file locking, and directory structures
- To explore file-system protection





File Concept

- Contiguous logical address space
- Types:
 - Data
 - numeric
 - character
 - binary
 - Program
- Contents defined by file's creator
 - Many types
 - Consider text file, source file, executable file





File Attributes

- Name only information kept in human-readable form
- Identifier unique tag (number) identifies file within file system
- Type needed for systems that support different types
- Location pointer to file location on device
- Size current file size
- Protection controls who can do reading, writing, executing
- **Time, date, and user identification** data for protection, security, and usage monitoring
- Information about files are kept in the directory structure, which is maintained on the disk
- Many variations, including extended file attributes such as file checksum
- Information kept in the directory structure

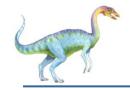




File info Window on Mac OS X



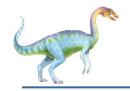




File Operations

- File is an abstract data type
- Create
- Write at write pointer location
- Read at read pointer location
- Reposition within file seek
- Delete
- Truncate
- Open (F_i) search the directory structure on disk for entry F_i , and move the content of entry to memory
- Close (F_i) move the content of entry F_i in memory to directory structure on disk

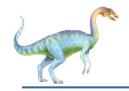




Open Files

- Several pieces of data are needed to manage open files:
 - Open-file table: tracks open files
 - File pointer: pointer to last read/write location, per process that has the file open
 - File-open count: counter of number of times a file is open – to allow removal of data from open-file table when last processes closes it
 - Disk location of the file: cache of data access information.
 - Access rights: per-process access mode information

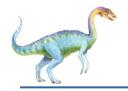




Open File Locking

- Provided by some operating systems and file systems
 - Similar to reader-writer locks
 - Shared lock similar to reader lock several processes can acquire concurrently
 - Exclusive lock similar to writer lock
- Mediates access to a file
- Mandatory or advisory:
 - Mandatory access is denied depending on locks held and requested
 - Advisory processes can find status of locks and decide what to do

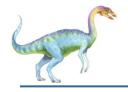




File Locking Example – Java API

```
import java.io.*;
import java.nio.channels.*;
public class LockingExample {
    public static final boolean EXCLUSIVE = false;
    public static final boolean SHARED = true;
    public static void main(String arsg[]) throws IOException {
            FileLock sharedLock = null;
            FileLock exclusiveLock = null;
           try {
                        RandomAccessFile raf = new RandomAccessFile("file.txt", "rw");
                       // get the channel for the file
                        FileChannel ch = raf.getChannel();
                       // this locks the first half of the file - exclusive
                        exclusiveLock = ch.lock(0, raf.length()/2, EXCLUSIVE);
                       /** Now modify the data . . . */
                       // release the lock
                        exclusiveLock.release();
```

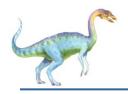




File Locking Example – Java API (Cont.)

```
// this locks the second half of the file - shared
          sharedLock = ch.lock(raf.length()/2+1, raf.length(),
                     SHARED);
          /** Now read the data . . . */
          // release the lock
          sharedLock.release();
} catch (java.io.IOException ioe) {
          System.err.println(ioe);
}finally {
          if (exclusiveLock != null)
          exclusiveLock.release();
          if (sharedLock != null)
          sharedLock.release();
```

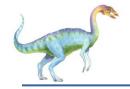




File Types – Name, Extension

file type	usual extension	function
executable	exe, com, bin or none	ready-to-run machine- language program
object	obj, o	compiled, machine language, not linked
source code	c, cc, java, pas, asm, a	source code in various languages
batch	bat, sh	commands to the command interpreter
text	txt, doc	textual data, documents
word processor	wp, tex, rtf, doc	various word-processor formats
library	lib, a, so, dll	libraries of routines for programmers
print or view	ps, pdf, jpg	ASCII or binary file in a format for printing or viewing
archive	arc, zip, tar	related files grouped into one file, sometimes com- pressed, for archiving or storage
multimedia	mpeg, mov, rm, mp3, avi	binary file containing audio or A/V information





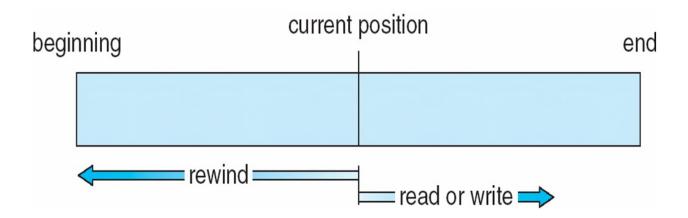
File Structure

- None sequence of words, bytes
- Simple record structure
 - Lines
 - Fixed length
 - Variable length
- Complex Structures
 - Formatted document
 - Relocatable load file
- Can simulate last two with first method by inserting appropriate control characters
- Who decides:
 - Operating system
 - Program





Sequential-access File







Access Methods

Sequential Access

```
read next
write next
reset
no read after last write
(rewrite)
```

■ **Direct Access** – file is fixed length logical records

```
read n
write n
position to n
read next
write next
rewrite n
```

n = relative block number

- Relative block numbers allow OS to decide where file should be placed
 - See allocation problem in Ch 12

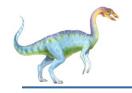




Simulation of Sequential Access on Direct-access File

sequential access	implementation for direct access	
reset	cp = 0;	
read next	read cp; cp = cp + 1;	
write next	write cp ; cp = cp + 1;	

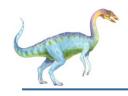




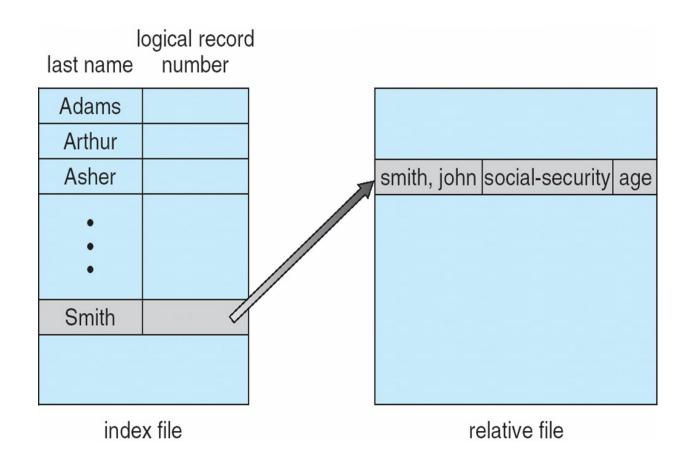
Other Access Methods

- Can be built on top of base methods
- General involve creation of an index for the file
- Keep index in memory for fast determination of location of data to be operated on (consider UPC code plus record of data about that item)
- If too large, index (in memory) of the index (on disk)
- IBM indexed sequential-access method (ISAM)
 - Small master index, points to disk blocks of secondary index
 - File kept sorted on a defined key
 - All done by the OS
- VMS operating system provides index and relative files as another example (see next slide)

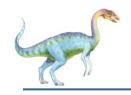




Example of Index and Relative Files

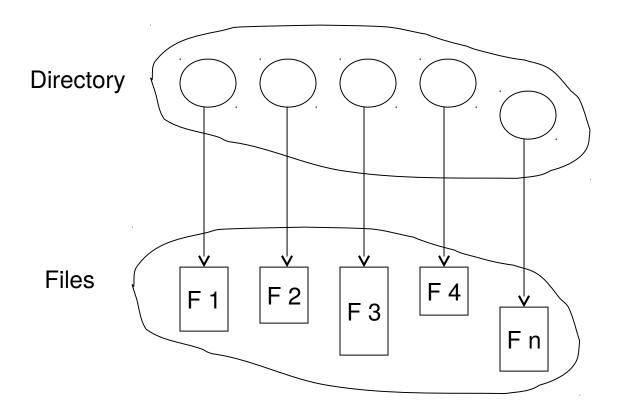




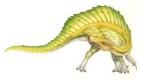


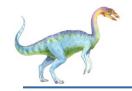
Directory Structure

A collection of nodes containing information about all files



Both the directory structure and the files reside on disk

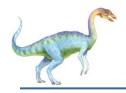




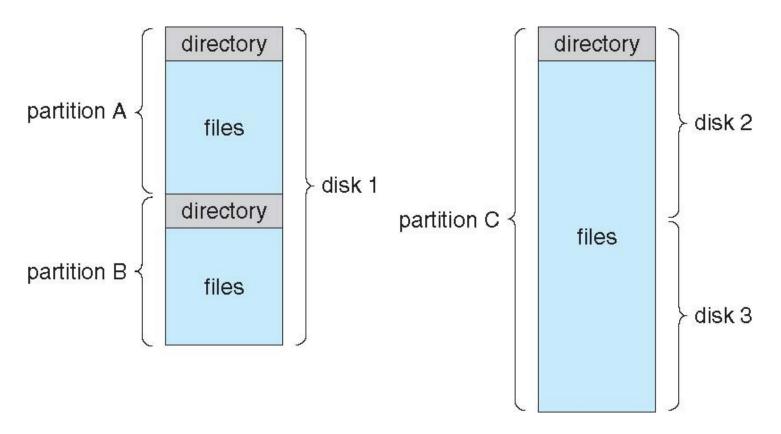
Disk Structure

- Disk can be subdivided into partitions
- Disks or partitions can be RAID protected against failure
- Disk or partition can be used raw without a file system, or formatted with a file system
- Partitions also known as minidisks, slices
- Entity containing file system known as a volume
- Each volume containing file system also tracks that file system's info in device directory or volume table of contents
- As well as general-purpose file systems there are many special-purpose file systems, frequently all within the same operating system or computer

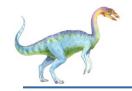




A Typical File-system Organization



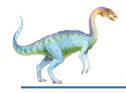




Types of File Systems

- We mostly talk of general-purpose file systems
- But systems frequently have may file systems, some general- and some special- purpose
- Consider Solaris has
 - tmpfs memory-based volatile FS for fast, temporary I/O
 - objfs interface into kernel memory to get kernel symbols for debugging
 - ctfs contract file system for managing daemons
 - lofs loopback file system allows one FS to be accessed in place of another
 - procfs kernel interface to process structures
 - ufs, zfs general purpose file systems

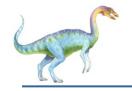




Operations Performed on Directory

- Search for a file
- Create a file
- Delete a file
- List a directory
- Rename a file
- Traverse the file system





Directory Organization

The directory is organized logically to obtain

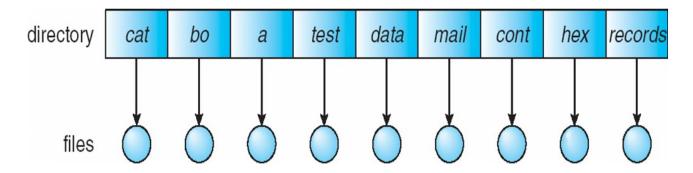
- Efficiency locating a file quickly
- Naming convenient to users
 - Two users can have same name for different files
 - The same file can have several different names
- Grouping logical grouping of files by properties, (e.g., all Java programs, all games, ...)





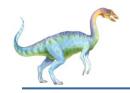
Single-Level Directory

A single directory for all users



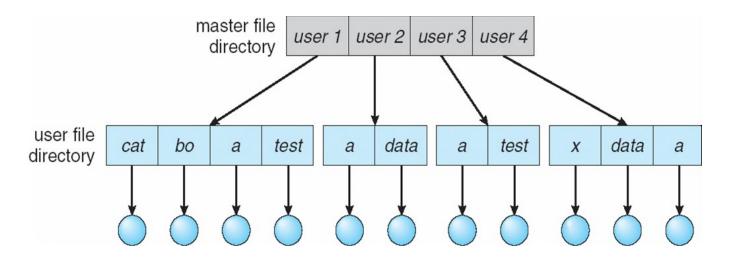
- Naming problem
- Grouping problem





Two-Level Directory

Separate directory for each user

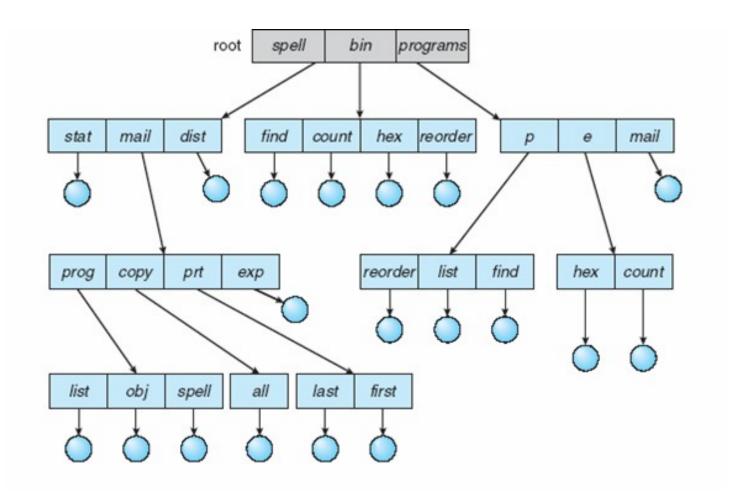


- Path name
- Can have the same file name for different user
- Efficient searching
- No grouping capability

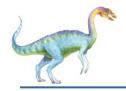




Tree-Structured Directories







Tree-Structured Directories (Cont.)

- Efficient searching
- Grouping Capability
- Current directory (working directory)
 - cd /spell/mail/prog
 - type list





Tree-Structured Directories (Cont)

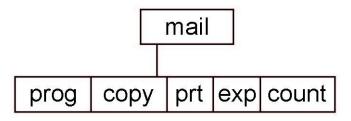
- Absolute or relative path name
- Creating a new file is done in current directory
- Delete a file

rm <file-name>

Creating a new subdirectory is done in current directory

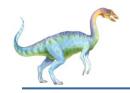
Example: if in current directory /mail

mkdir count



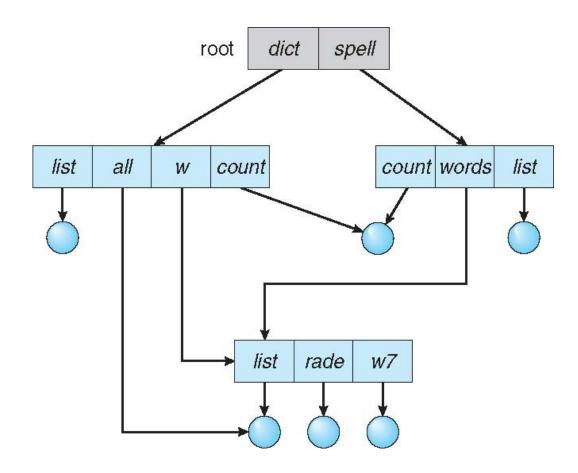
Deleting "mail" ⇒ deleting the entire subtree rooted by "mail"



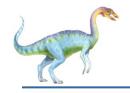


Acyclic-Graph Directories

Have shared subdirectories and files







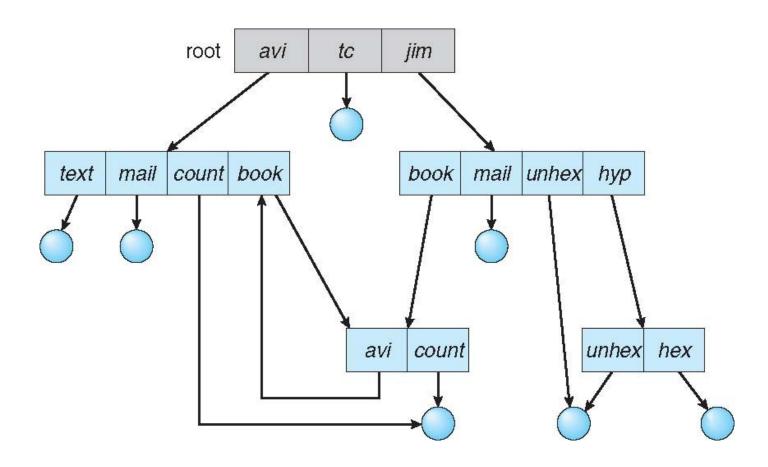
Acyclic-Graph Directories (Cont.)

- Two different names (aliasing)
- If dict deletes list ⇒ dangling pointer Solutions:
 - Backpointers, so we can delete all pointers
 Variable size records a problem
 - Backpointers using a daisy chain organization
 - Entry-hold-count solution
- New directory entry type
 - Link another name (pointer) to an existing file
 - Resolve the link follow pointer to locate the file

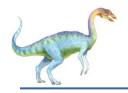




General Graph Directory



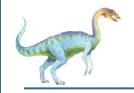




General Graph Directory (Cont.)

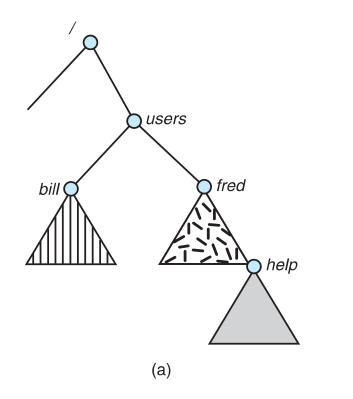
- How do we guarantee no cycles?
 - Allow only links to file not subdirectories
 - Garbage collection
 - Every time a new link is added use a cycle detection algorithm to determine whether it is OK

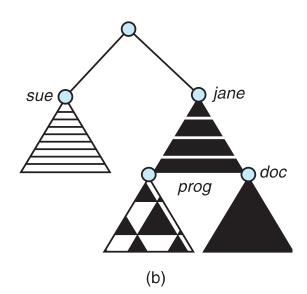




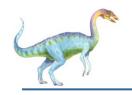
File System Mounting

- A file system must be mounted before it can be accessed
- A unmounted file system (i.e., Fig. 11-11(b)) is mounted at a mount point

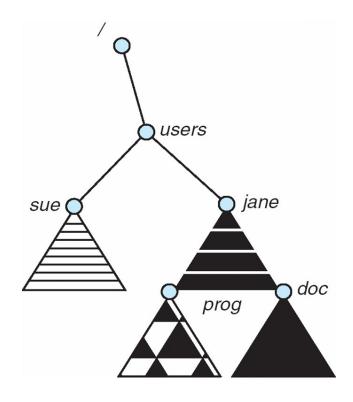




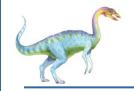




Mount Point







File Sharing

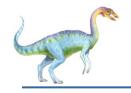
- Sharing of files on multi-user systems is desirable
- Sharing may be done through a protection scheme
- On distributed systems, files may be shared across a network
- Network File System (NFS) is a common distributed file-sharing method
- If multi-user system
 - User IDs identify users, allowing permissions and protections to be per-user
 Group IDs allow users to be in groups, permitting group access rights
 - Owner of a file / directory
 - Group of a file / directory





File Sharing – Remote File Systems

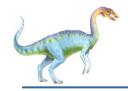
- Uses networking to allow file system access between systems
 - Manually via programs like FTP
 - Automatically, seamlessly using distributed file systems
 - Semi automatically via the world wide web
- Client-server model allows clients to mount remote file systems from servers
 - Server can serve multiple clients
 - Client and user-on-client identification is insecure or complicated
 - NFS is standard UNIX client-server file sharing protocol
 - CIFS is standard Windows protocol
 - Standard operating system file calls are translated into remote calls
- Distributed Information Systems (distributed naming services) such as LDAP, DNS, NIS, Active Directory implement unified access to information needed for remote computing



File Sharing – Failure Modes

- All file systems have failure modes
 - For example corruption of directory structures or other nonuser data, called metadata
- Remote file systems add new failure modes, due to network failure, server failure
- Recovery from failure can involve state information about status of each remote request
- Stateless protocols such as NFS v3 include all information in each request, allowing easy recovery but less security

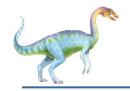




File Sharing – Consistency Semantics

- Specify how multiple users are to access a shared file simultaneously
 - Similar to Ch 5 process synchronization algorithms
 - Tend to be less complex due to disk I/O and network latency (for remote file systems
 - Andrew File System (AFS) implemented complex remote file sharing semantics
 - Unix file system (UFS) implements:
 - Writes to an open file visible immediately to other users of the same open file
 - Sharing file pointer to allow multiple users to read and write concurrently
 - AFS has session semantics
 - Writes only visible to sessions starting after the file is closed

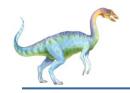




Protection

- File owner/creator should be able to control:
 - what can be done
 - by whom
- Types of access
 - Read
 - Write
 - Execute
 - Append
 - Delete
 - List



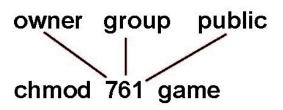


Access Lists and Groups

- Mode of access: read, write, execute
- Three classes of users on Unix / Linux

			RWX
a) owner access	7	\Rightarrow	111
,			RWX
b) group access	6	\Rightarrow	110
			RWX
c) public access	1	\Rightarrow	001

- Ask manager to create a group (unique name), say G, and add some users to the group.
- For a particular file (say *game*) or subdirectory, define an appropriate access.



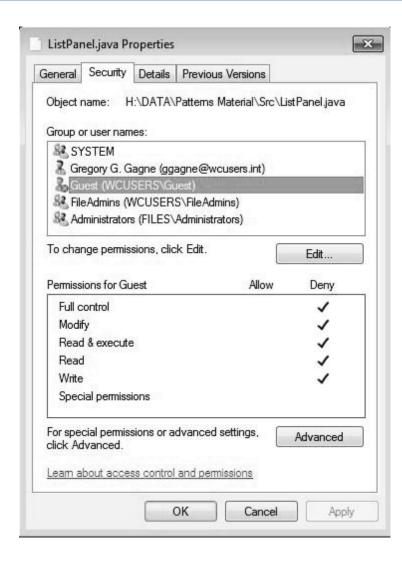
Attach a group to a file

chgrp G game





Windows 7 Access-Control List Management







A Sample UNIX Directory Listing

-rw-rw-r	1 pbg	staff	31200	Sep 3 08:30	intro.ps
drwx	5 pbg	staff	512	Jul 8 09.33	private/
drwxrwxr-x	2 pbg	staff	512	Jul 8 09:35	doc/
drwxrwx	2 pbg	student	512	Aug 3 14:13	student-proj/
-rw-rr	1 pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1 pbg	staff	20471	Feb 24 2003	program
drwxxx	4 pbg	faculty	512	Jul 31 10:31	lib/
drwx	3 pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3 pbg	staff	512	Jul 8 09:35	test/



End of Chapter 13

